



Nonlinear pattern matching in rule-based modeling languages

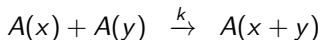
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Rule-based modeling languages

Stochastic graph rewriting: Kappa, BNGL

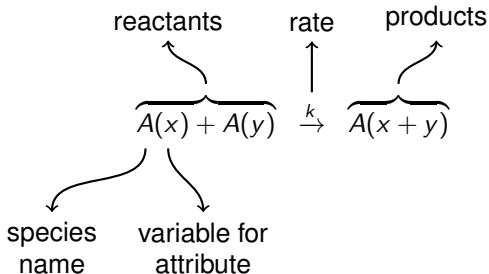
Stochastic term rewriting: (C)SMMR, ML-Rules, React(C), Chromar



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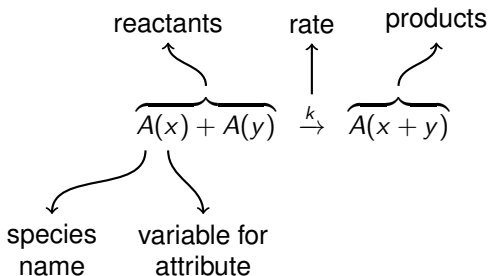
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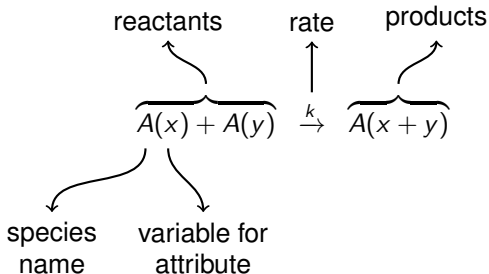


Match in solution
 $\{A(1), A(2), \dots\}$

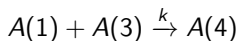
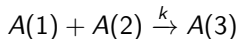
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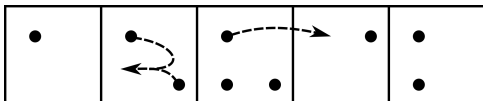


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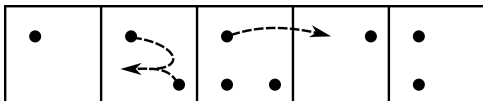
Overview

- Rules use patterns on their left side, which are matched to the current solution to obtain reactions.
- **Nonlinear patterns** are patterns in which variables may occur multiple times.
- In this talk:
 - Nonlinear patterns are useful for expressing relations between reactants (e.g., **spatial relations**).
 - State-of-the-art languages only support linear patterns.
 - Supporting nonlinear patterns allows more efficient pattern matching.

The case for nonlinear patterns



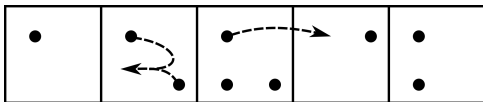
The case for nonlinear patterns



$$A(x) + A(x) \rightarrow \emptyset$$

$$A(x) \rightarrow A(x + 1)$$

The case for nonlinear patterns



$$A(x) + A(x) \rightarrow \emptyset$$

$$A(x) \rightarrow A(x + 1)$$

$$C(x)[A + A + r] \rightarrow C(x)[r]$$

$$C(x)[A + r_1] + C(x + 1)[r_2] \rightarrow C(x)[r_1] + C(x + 1)[A + r_2]$$

State of the art

Expressing patterns

State-of-the-art languages only allow linear patterns, as linear pattern matching is simpler than nonlinear pattern matching.

In these languages, nonlinear patterns must be **linearized** to express them, resulting in rules with **constraints**.

$$A(x) + A(x) \longrightarrow \emptyset$$

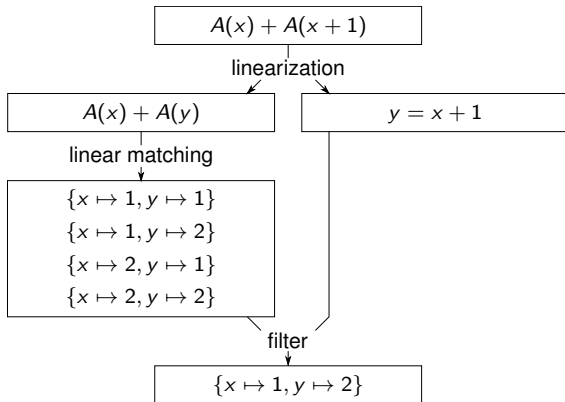
$$A(x) + A(y) \xrightarrow{x=y} \emptyset$$

$$C(x)[A + r_1] + C(x + 1)[r_2] \longrightarrow C(x)[r_1] + C(x + 1)[A + r_2]$$

$$C(x)[A + r_1] + C(y)[r_2] \xrightarrow{y=x+1} C(x)[r_1] + C(y)[A + r_2]$$

State of the art

Matching $A(x) + A(x + 1)$ in the solution $\{A(1), \dots, A(n)\}$ with $n = 2$





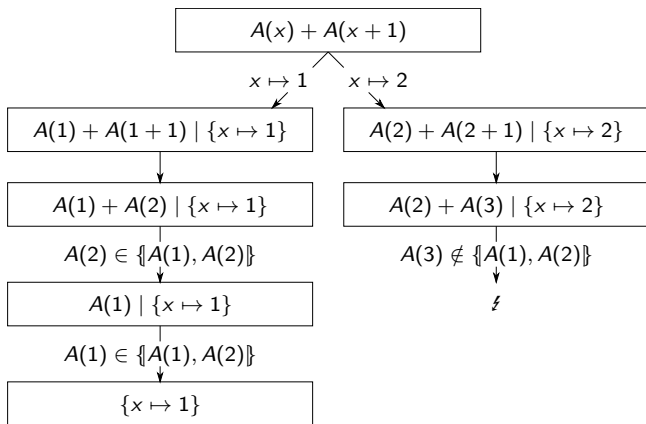
State of the art

Problem

- In a solution $\{A(1), \dots, A(n)\}$ we get $n - 1$ eventual matches, but n^2 intermediate results.
- Most intermediate results are thrown away.
- The relation between different occurrences of the same variable is not available during pattern matching.

Inline substitution

Matching $A(x) + A(x + 1)$ in the solution $\{A(1), \dots, A(n)\}$ with $n = 2$





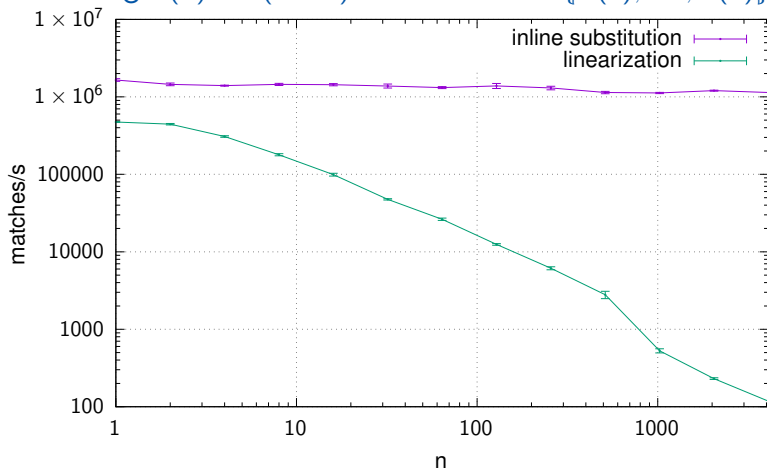
Inline substitution

Idea

Try all possible values for a variable

- The chosen value is substituted for **all occurrences** of the variable.
- ⇒ We can **evaluate expressions**.
- ⇒ We can **quickly** check whether the reactants exist.
- ⇒ We can prune unsuccessful branches early.

Matching $A(x) + A(x + 1)$ in the solution $\{A(1), \dots, A(n)\}$





Context

- Nonlinear patterns are unpopular in **term rewriting** and functional programming.
- Logic programming languages (e.g., Prolog) allow nonlinear patterns, but have limited support for expressions.
- **Functional logic programming** languages (e.g., Curry) are a current research topic.
- We need to count match multiplicities for mass action kinetics.

Conclusion and open questions

- Usually, a **more expressive** language means **less efficient** execution.
- However, nonlinear patterns contain valuable information about the **modeler intent**, enabling faster algorithms.
- Can we relate term rewriting and graph rewriting?
- Is pattern matching decisive for simulation performance?
- What is easier to read and write?

$$A(x) + A(x + 1) \longrightarrow \emptyset \qquad A(x) + A(y) \xrightarrow{y=x+1} \emptyset$$